

9-8-03 CSE571 notes

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- $AnsProlog^{-not} : a_0 \leftarrow a_1 \dots a_m.$
- $AnsProlog : a_0 \leftarrow a_1 \dots a_m, nota_{m+1} \dots nota_n.$
- Minimal Model - there is no smaller model, there is no proper subset
- Least Model - there is no smaller model or model of the same size
- A model is an interpretation that satisfies all rules
- A rule is satisfied if the RHS is true, or we don't need to prove the LHS.
- Examples:
 - Ex1.
 $a \leftarrow b.$
 $HB = \{a, b\}$
 $HI = \{\emptyset, \{a\}, \{b\}, \{a, b\}\}$
 $HMs = \{\emptyset, \{a\}, \{a, b\}\}$
 $MMs = \{\emptyset\}$
 $LM = \{\emptyset\}$
 - Ex2.
 $a \leftarrow notb.$
 $b \leftarrow nota.$
 $HB = \{a, b\}$
 $HI = \{\emptyset, \{a\}, \{b\}, \{a, b\}\}$
 $HMs = \{\{a\}, \{b\}, \{a, b\}\}$
 $MMs = \{\{a\}, \{b\}\}$
 $LM = none$
 - Ex3.
 $a \leftarrow notb.$
 $b \leftarrow nota.$
 $b \leftarrow c.$
 $c \leftarrow b.$
 $HB = \{a, b, c\}$
 $HI = \{\emptyset, \{a\}, \{b\}, \{c\}, \{a, b\}, \{a, c\}, \{b, c\}, \{a, b, c\}\}$
 $HMs = \{\{a\}, \{b, c\}\}$
 $MMs = \{\{a\}, \{b, c\}\}$
 $LM = \{a\}$

– Ex4.

$a \leftarrow notb.$

$HB = \{a, b\}$

$HI = \{\emptyset, \{a\}, \{b\}, \{a, b\}\}$

$HMs = \{\{a\}, \{b\}\}$

$MMs = \{\{a\}, \{b\}\}$

$LM = none$

Here even though $\{b\}$ is a minimal model, it is not stable.

- For $AnsProlog^{-not}$, we can define the notion of stable models, which is equivalent to the following terms:

- answer set
- minimal models
- least model
- least fixpoint of T_{Π}^0 operator

- Ex5.

$a \leftarrow b.$

$b \leftarrow c, d.$

$d \leftarrow c.$

$c \leftarrow e.$

$f \leftarrow g$

$e \leftarrow .$

$HB = \{a, b, c, d, e, f, g\}$

$HI = \{\emptyset, \{a\}, \{b\}, \{c\}, \{a, b\}, \{a, c\}, \{b, c\}, \{a, b, c\}, \dots\}$

- The fixpoint is found by applying all of the applicable rules at each iteration until two iterations are identical.

- The fixpoint procedure for Ex5. is as follows:

\emptyset

e

e, c

e, c, d

e, c, d, b

e, c, d, b, a

e, c, d, b, a

- If nothing can be derived from \emptyset , then \emptyset is the least model.
- Formally $T_{\Pi}^0(I) = \{L_0 \in HB_{\Pi} \text{ s.t. } L_0 \leftarrow L_1 \dots L_m \text{ is a rule in } \Pi \text{ and } \{L_1 \dots L_m\} \subseteq I\}$
- For $AnsProlog$, $T_{\Pi}^1 = \{L_0 \text{ s.t. } \exists \text{ a rule with } L_0 \text{ in the head and the body is true w.r.t. } I\}$
- Why finding fixed points is not reasonable in $AnsProlog$: Consider Ex2, starting with \emptyset , you derive $\{a, b\}$, from which you derive \emptyset . There is no fix point because it will cycle forever. The next best idea is to take an interpretation and check if it is a stable model by converting it to a $AnsProlog^{-not}$ program.

- Converting AnsProlog to *AnsProlog*^{-not}: Given an AnsProlog program Π and an interpretation I , Π^I is an *AnsProlog*^{-not} program obtained from Π by:
 1. (i) Removing rules of the form $l_0 \leftarrow l_1 \dots l_m, \text{not} l_{m+1} \dots \text{not} l_n$ from Π if for some $j : m + 1 \leq j \leq n, l_j \in I$.
 2. (ii) Removing, for all x , the parts of rules stating 'not x '.
- I is an answer set of an AnsProlog program Π if I is the least model of the *AnsProlog*^{-not} program Π^I .

- Ex6.

$\Pi = a \leftarrow \text{not} b.$

$b \leftarrow \text{not} a.$

$HB = \{a, b\}$

$HI = \{\emptyset, \{a\}, \{b\}, \{a, b\}\} = \{I_1, I_2, I_3, I_4\}$

Find the answer sets for each Π^{I_j} :

Initially $\forall_j \Pi = \Pi^{I_j}$

After rule (i)

$\Pi^{I_1} = \Pi$

$\Pi^{I_2} = a \leftarrow \text{not} b.$

$\Pi^{I_3} = b \leftarrow \text{not} a.$

$\Pi^{I_4} = \emptyset$

After rule (ii)

$\Pi^{I_1} = a \leftarrow .$

$b \leftarrow .$

$\Pi^{I_2} = a \leftarrow .$

$\Pi^{I_3} = b \leftarrow .$

$\Pi^{I_4} = \emptyset$

The least models are:

$lm^{I_1} = \{a, b\}$

$lm^{I_2} = \{a\}$

$lm^{I_3} = \{b\}$

$lm^{I_4} = \emptyset$

Of these, only lm^{I_2} and lm^{I_3} are answer sets because the least model of the *AnsProlog*^{-not} program is equal to the interpretation of the AnsProlog program.