

CSE355 Fall 2016–Recitation Quiz 14 (Solutions)

Name: _____ ASU ID: _____

Monday 9:40AM Monday 10:45AM Tuesday 7:30AM Tuesday 4:35PM Wednesday 9:40AM

Wednesday 10:45AM Thursday 7:30AM Thursday 4:35PM Friday 9:40AM Friday 10:45AM

1. Show that $INFINITE_{TM} = \{\langle M \rangle \mid M \text{ is a TM and } L(M) \text{ is infinite}\}$ is undecidable by a reduction from A_{TM} .

Suppose $INFINITE_{TM}$ were decided by I . Then we can make a decider D for A_{TM} as follows:

$D =$ “On input $\langle M, w \rangle$ where M is a TM:

- (a) Construct a TM $M' =$ “On input x :
 - i. If $x = \epsilon$, accept.
 - ii. If $x \neq \epsilon$, run M on w and accept x only if M accepts w .”
- (b) Run I on $\langle M' \rangle$.
- (c) If I accepts, accept; otherwise, reject.”

If M accepts w , then $L(M') = \Sigma^*$, which is an infinite language; if M does not accept w , then $L(M') = \{\epsilon\}$, which is a finite language. Therefore, since A_{TM} is undecidable, then so is $INFINITE_{TM}$.

In fact, this actually shows something more important: given *any* language $L \subseteq \Sigma^*$, it is undecidable to check that an arbitrary TM M has $L(M) = L$.

2. Answer True or False and give a brief explanation as to your answer:

- (a) *Undecidable* languages are closed under complement.

True. If they were not closed under complement, then some undecidable language L has \bar{L} being decidable. But then since decidable languages are closed under complement, we have L being decidable, a contradiction.

- (b) If $A \leq_m B$ and B is not decidable, then A is not decidable.

False. We cannot say anything if B is not decidable and $A \leq_m B$. However, we *can* say something if instead of B being undecidable, A was undecidable. So therefore, the reduction is “in the wrong direction.”

- (c) If $A \leq_m \bar{A}$ and A is recognizable, then A is decidable.

True. We also then have $\bar{A} \leq_m A$ with the same reduction. Since A is recognizable, then because of the result of $A \leq_m B$ with A being recognizable implies B recognizable, then \bar{A} is recognizable. But we know that if a language and its complement are both recognizable, then it is decidable.

TA use only! Quiz has been recorded: