

# Learning Congestion Patterns in Optimal Power Flow Problems

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# Learning Congestion Patterns in Optimal Power Flow Problems



with Sidhant Misra (LANL), Yeesian Ng (MIT), Sean Simpson (UW Madison), and Dan Molzahn (Georgia Tech)



https://electrical-engineering-portal.com/wp-content/uploads/2018/02/power-transmission-control-center.jpg



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min  $\sum_{i \in G} (c_{2,i} p_{G,i}^2 + c_{1,i} p_{G,i})$ s.t.  $p_{G(i)} - p_{D(i)} = \sum_{j \in \delta i} b_{ij} (\theta_i - \theta_j)$  ,  $p_{G,g}^{min} \leq p_{G,g} \leq p_{G,g}^{max}$ ,  $g \in \mathcal{G}$  $-f_{ij}^{max} \le b_{ij} (\theta_i - \theta_j) \le f_{ij}^{max}, j \in \mathcal{L}$ 

#### "Power system operators solve power flow in their heads"



min  $\sum_{i \in G} (c_{2,i} p_{G,i}^2 + c_{1,i} p_{G,i})$ s.t.  $p_{G(i)} - p_{D(i)} = \sum_{j \in \delta i} b_{ij} (\theta_i - \theta_j)$  ,  $p_{G,g}^{min} \leq p_{G,g} \leq p_{G,g}^{max}$ ,  $g \in \mathcal{G}$  $-f_{ij}^{max} \leq b_{ij} (\theta_i - \theta_j) \leq f_{ij}^{max}, j \in \mathcal{L}$ 

IDEA: Can we use *machine learning* to identify the optimal dispatch?

https://electrical-engineering-portal.com/wp-content/uploads/2018/02/power-transmission-control-center.jpg

Goal: Low cost operation, while enforcing technical limits

min  $\sum_{i \in G} (c_{2,i} p_{G,i}^2 + c_{1,i} p_{G,i})$  minimize generation cost

s.t.

 $p_{G(i)} - p_{D(i)} = \sum_{j \in \delta i} b_{ij} (\theta_i - \theta_j)$ , power flow balance

 $p_{G,g}^{min} \le p_{G,g} \le p_{G,g}^{max}$ ,  $g \in \mathcal{G}$  generation constraints

 $-f_{ij}^{max} \leq b_{ij} (\theta_i - \theta_j) \leq f_{ij}^{max}, j \in \mathcal{L}$ 

transmission constraints

Goal: Low cost operation, while enforcing technical limits

min 
$$\sum_{i \in \mathcal{G}} (c_{2,i} p_{G,i}^2 + c_{1,i} p_{G,i})$$
 minimize generation cost

s.t.

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 $p_{G,g}^{min} \le p_{G,g} \le p_{G,g}^{max}$ ,  $g \in \mathcal{G}$  generation constraints

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Could include much more details (security constraints, generator variables on/off variables, non-linear AC power flow, ...)

System conditions (load, renewables) are constantly changing

min 
$$\sum_{i \in G} (c_{2,i} p_{G,i}^2 + c_{1,i} p_{G,i})$$
 minimize generation cost

s.t.

 $p_{G(i)} - \mathbf{p}_{\mathbf{D}(i)} = \sum_{j \in \delta i} b_{ij} (\theta_i - \theta_j)$ , power flow balance

 $p_{G,g}^{min} \le p_{G,g} \le p_{G,g}^{max}$ ,  $g \in \mathcal{G}$  generation constraints

 $-f_{ij}^{max} \le b_{ij} (\theta_i - \theta_j) \le f_{ij}^{max}, j \in \mathcal{L}$  transmission constraints

Could include much more details (security constraints, generator variables on/off variables, non-linear AC power flow, ...)

System conditions (load, renewables) are constantly changing

$$\min \sum_{i \in \mathcal{G}} (c_{2,i} p_{G,i}^2 + c_{1,i} p_{G,i})$$
minimize generation cost
$$s.t. \\ p_{G(i)} - p_{D(i)} = \sum_{j \in \delta i} b_{ij}(\theta_i - \theta_j),$$
power flow balance
$$p_{G,g}^{min} \leq p_{G,g} \leq p_{G,g}^{max}, g \in \mathcal{G}$$
generation constraints
$$-f_{ij}^{max} \leq b_{ij}(\theta_i - \theta_j) \leq f_{ij}^{max}, j \in \mathcal{L}$$
transmission constraints
$$FResolve this problem every 5-15 \min for varying load profile p_D$$

Could include much more details (security constraints, generator variables on/off variables, non-linear AC power flow, ...)

```
OPF at T_1 with load p_D^{(1)}

min \sum_{i \in \mathcal{G}} (c_{2,i} p_{G,i}^2 + c_{1,i} p_{G,i})

s.t.

p_{G(i)} - p_{D(i)} = \sum_{j \in \delta i} b_{ij} (\theta_i - \theta_j),

p_{G,g}^{min} \leq p_{G,g} \leq p_{G,g}^{max}, g \in \mathcal{G}

-f_{ij}^{max} \leq b_{ij} (\theta_i - \theta_j) \leq f_{ij}^{max}, j \in \mathcal{L}
```

Optimal solution  $p_G^*\left(p_D^{(1)}\right)$ 

OPF at  $T_1$  with load  $p_D^{(1)}$ min  $\sum_{i \in G} (c_{2,i} p_{G,i}^2 + c_{1,i} p_{G,i})$ s.t.  $p_{G(i)} - p_{D(i)} = \sum_{j \in \delta i} b_{ij} (\theta_i - \theta_j),$   $p_{G,g}^{min} \leq p_{G,g} \leq p_{G,g}^{max}, g \in G$  $-f_{ij}^{max} \leq b_{ij} (\theta_i - \theta_j) \leq f_{ij}^{max}, j \in \mathcal{L}$  OPF at  $T_2$  with load  $p_D^{(2)}$ 

 $\min \sum_{i \in \mathcal{G}} (c_{2,i} p_{G,i}^2 + c_{1,i} p_{G,i})$ s.t.  $p_{G(i)} - \mathbf{p}_{\mathbf{D}(i)} = \sum_{j \in \delta i} b_{ij} (\theta_i - \theta_j),$   $p_{G,g}^{min} \le p_{G,g} \le p_{G,g}^{max}, g \in \mathcal{G}$   $-f_{ij}^{max} \le b_{ij} (\theta_i - \theta_j) \le f_{ij}^{max}, j \in \mathcal{L}$ 

Optimal solution  $p_{G}^{*}\left(p_{D}^{(1)}\right)$ 

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OPF at  $T_1$  with load  $p_D^{(1)}$ min  $\sum_{i \in \mathcal{G}} (c_{2,i} p_{G,i}^2 + c_{1,i} p_{G,i})$ s.t.  $p_{G(i)} - p_{D(i)} = \sum_{j \in \delta i} b_{ij} (\theta_i - \theta_j),$   $p_{G,g}^{min} \leq p_{G,g} \leq p_{G,g}^{max}, g \in \mathcal{G}$  $-f_{ij}^{max} \leq b_{ij} (\theta_i - \theta_j) \leq f_{ij}^{max}, j \in \mathcal{L}$  OPF at  $T_2$  with load  $p_p^{(2)}$ 

 $\min \sum_{i \in \mathcal{G}} (c_{2,i} p_{G,i}^2 + c_{1,i} p_{G,i})$ s.t.  $p_{G(i)} - \mathbf{p}_{D(i)} = \sum_{j \in \delta i} b_{ij} (\theta_i - \theta_j),$   $p_{G,g}^{min} \le p_{G,g} \le p_{G,g}^{max}, g \in \mathcal{G}$   $-f_{ij}^{max} \le b_{ij} (\theta_i - \theta_j) \le f_{ij}^{max}, j \in \mathcal{L}$  OPF at  $T_3$  with load  $p_D^{(3)}$ min  $\sum_{i \in \mathcal{G}} (c_{2,i} p_{G,i}^2 + c_{1,i} p_{G,i})$ s.t.  $p_{G(i)} - p_{D(i)} = \sum_{j \in \delta i} b_{ij} (\theta_i - \theta_j),$   $p_{G,g}^{min} \leq p_{G,g} \leq p_{G,g}^{max}, g \in \mathcal{G}$  $-f_{ij}^{max} \leq b_{ij} (\theta_i - \theta_j) \leq f_{ij}^{max}, j \in \mathcal{L}$ 

Optimal solution  $p_G^*\left(p_D^{(1)}\right)$ 

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Optimal solution  $p_{G}^{*}\left(p_{D}^{(1)}\right)$ 

Optimal solution  $p_{G}^{*}\left(p_{D}^{(2)}\right)$ 

Optimal solution  $p_G^*\left(p_D^{(3)}\right)$ 



#### Can we use solutions for previous load profiles $p_D$

to learn the new solution  $p_G^*$  ?

First attempt:

Train a neural net!





• This didn't work well...

(though I will admit that we gave up quite fast)



- This didn't work well...
  - Hard to satisfy safety constraints!



- This didn't work well...
  - Hard to satisfy safety constraints!
  - Projection back onto feasible space cause suboptimality...



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  - Challenging: High-dimensional input  $\rightarrow$  High dimensional output



- This didn't work well...
  - Hard to satisfy safety constraints!
  - Projection back onto feasible space cause suboptimality...
  - Challenging: High-dimensional input  $\rightarrow$  High dimensional output
- This can work well under some circumstances

Wide enough and deep enough [Karg and Lucia, 2018]



https://electrical-engineering-portal.com/wp-content/uploads/2018/02/power-transmission-control-center.jpg

 $\min_{p_G,\theta} \quad \sum_{i\in\mathcal{G}} c_{1,i} \, p_{G,i}$ s.t.  $p_{G(i)} - p_{D(i)} = \sum_{i \in \delta i} b_{ij} (\theta_i - \theta_j),$  $p_{G,g}^{min} \leq p_{G,g} \leq p_{G,g}^{max}, g \in \mathcal{G}$  $-f_{ij}^{max} \leq b_{ij} (\theta_i - \theta_j) \leq f_{ij}^{max}, j \in \mathcal{L}$ 

#### "Power system operators solve power flow in their heads"

#### How is this possible?

Limited number of **operational patterns** 

- Set of congested lines
- Set of generators at their limits

/	
	$\min_{p_G,\theta} \sum_{i\in \mathcal{G}} c_{1,i} p_{G,i}$
	s.t. $\sum_{n=1}^{\infty} h(0, 0)$
	$p_{G(i)} - p_{D(i)} = \sum_{j \in \delta i} p_{ij} (\theta_i - \theta_j),$
	$p_{G,g}^{min} \leq p_{G,g} \leq p_{G,g}^{max}, \ g \in \mathcal{G}$
T	$-f_{ij}^{max} \le b_{ij} (\theta_i - \theta_j) \le f_{ij}^{max}, j \in \mathcal{L}$

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This corresponds to the active constraints of an OPF problem!

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Operational pattern = optimal active set  $A^*$ 

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T	$-f_{ij}^{max} \le b_{ij} (\theta_i - \theta_j) \le f_{ij}^{max}, j \in \mathcal{L}$

Shorthand version of the OPF:

min  $\sum_{i \in \mathcal{G}} (c_{2,i} p_{G,i}^2 + c_{1,i} p_{G,i})$ 

s.t.  $Ap_G \leq b + Cp_D$ 

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This corresponds to the active constraints of an OPF problem!

Operational pattern = optimal active set  $A^*$ 

	$\min_{p_G,\theta} \sum_{i\in \mathcal{G}} c_{1,i} p_{G,i}$
	s.t. $p_{1} = p_{2} = \sum_{i=1}^{n} b_{i}(\theta_{i} - \theta_{i})$
	$p_{G(i)}^{min} < p_{C,a} < p_{C,a}^{max},  a \in G$
7-	$-f_{ii}^{max} \le b_{ii}(\theta_i - \theta_i) \le f_{ii}^{max}, j \in \mathcal{L}$
	·

Shorthand version of the OPF:

min  $\sum_{i \in \mathcal{G}} (c_{2,i} p_{G,i}^2 + c_{1,i} p_{G,i})$ 

s.t.  $Ap_G \leq b + Cp_D$ 

Active constaints are satisfied with equality  $A_{act}p_G = b_{act} + C_{act}p_D$ 

#### How is this possible?

Limited number of operational patterns

- Set of congested lines
- Set of generators at their limits

Each optimal operational pattern  $\mathcal{A}^*$  has a corresponding **optimal control policy** 

 $\boldsymbol{p}_{\boldsymbol{G}}^* = A_{act}^{-1}(b_{act} + C_{act}\boldsymbol{p}_{\boldsymbol{D}})$ 

$\min_{p_G,\theta} \sum_{i\in \mathcal{G}} c_{1,i} p_{G,i}$
s.t. $p_{G(i)} - p_{D(i)} = \sum_{j \in \delta i} b_{ij} (\theta_i - \theta_j),$
$p_{G,g}^{min} \leq p_{G,g} \leq p_{G,g}^{max}, \ g \in \mathcal{G}$
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Shorthand version of the OPF:

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Each optimal operational pattern *A*\* has a corresponding *optimal control policy* 

 $\boldsymbol{p}_{\boldsymbol{G}}^* = A_{act}^{-1}(b_{act} + C_{act}\boldsymbol{p}_{\boldsymbol{D}})$ 

Optimal dispatch is a *linear function* of the load profile  $p_D$  for *small*\* changes!

\*Small changes = same active constraints at optimal solution



#### How is this possible?

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Each optimal operational pattern *A*\* has a corresponding *optimal control policy* 

 $\boldsymbol{p}_{\boldsymbol{G}}^* = A_{act}^{-1}(b_{act} + C_{act}\boldsymbol{p}_{\boldsymbol{D}})$ 

Multiparametric programming, explicit MPC... [Ji, Thomas and Tong '16], [Geng and Xie '17]

 $\min_{p_G,\theta} \sum_{i\in\mathcal{G}} c_{1,i} p_{G,i}$ s.t.  $p_{G(i)} - p_{D(i)} = \sum_{j \in \delta i} b_{ij} (\theta_i - \theta_j)$  ,  $p_{G,g}^{min} \leq p_{G,g} \leq p_{G,g}^{max}, g \in \mathcal{G}$  $-f_{ij}^{max} \le b_{ij} (\theta_i - \theta_j) \le f_{ij}^{max}, j \in \mathcal{L}$ (b) 30 20  $\Delta P_{w2}$ -10 [Vrakopoulou and Hiskens, 2017] -20 -15 -10 Δ P

Multi-parametric programming is inconvenient and computationally expensive for large systems 🛞

The number of possible active sets is exponential in problem size 😕

Multiparametric programming, explicit MPC... [Ji, Thomas and Tong '16], [Geng and Xie '17]



Multi-parametric programming is inconvenient and computationally expensive for large systems 🛞

The number of possible active sets is exponential in problem size 🐵

... but maybe only a few matters in practice? ©

Can we learn the practically relevant active sets?



# Think again!



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 $\min_{p_G,\theta} \sum_{i\in \mathcal{G}} c_{1,i} p_{G,i}$ s.t.  $p_{G(i)} - p_{D(i)} = \sum_{j \in \delta i} b_{ij} (\theta_i - \theta_j)$  ,  $p_{G,g}^{min} \leq p_{G,g} \leq p_{G,g}^{max}, g \in \mathcal{G}$  $-f_{ij}^{max} \leq b_{ij} (\theta_i - \theta_j) \leq f_{ij}^{max}, j \in \mathcal{L}$ 

IDEA: Can we use *statistical learning* to identify optimal active sets and corresponding optimal policies?

### Learning the Patterns

Learning optimal active sets through sampling

 $\begin{array}{ccc} \text{load} & & \text{Optimal} \\ \text{profile 1} \longrightarrow & \text{Power} \\ p_D & & \text{Flow} \end{array} \xrightarrow{p_G^*} \text{Optimal dispatch} \\ \mathcal{A}^* \text{ Optimal active set 1} \end{array}$ 

$$\begin{split} \min_{p_{G},\theta} & \sum_{i \in \mathcal{G}} c_{1,i} p_{G,i} \\ \text{s.t.} & p_{G(i)} - \boldsymbol{p}_{\boldsymbol{D}(i)} = \sum_{j \in \delta i} b_{ij} (\theta_{i} - \theta_{j}), \\ & p_{G,g}^{min} \leq p_{G,g} \leq p_{G,g}^{max}, \\ & -f_{ij}^{max} \leq b_{ij} (\theta_{i} - \theta_{j}) \leq f_{ij}^{max}, \\ & i \in N, g \in \mathcal{G}, ij \in \mathcal{L} \end{split}$$

### Learning the Patterns

Learning optimal active sets through sampling





### Learning the Patterns

Learning optimal active sets through sampling




Learning optimal active sets through sampling

Flow

. . .

 $p_D$ 



 $\min_{p_{G},\theta} \sum_{i \in \mathcal{G}} c_{1,i} p_{G,i}$ s.t.  $p_{G(i)} - \mathbf{p}_{\mathbf{D}(i)} = \sum_{j \in \delta i} b_{ij} (\theta_{i} - \theta_{j}),$  $p_{G,g}^{min} \leq p_{G,g} \leq p_{G,g}^{max},$  $-f_{ij}^{max} \leq b_{ij} (\theta_{i} - \theta_{j}) \leq f_{ij}^{max},$  $i \in N, g \in \mathcal{G}, ij \in \mathcal{L}$ 

 $\mathcal{A}^*$  Optimal active set 3

Learning optimal active sets through sampling

. . .



 $\min_{p_G,\theta} \sum_{i\in\mathcal{G}} c_{1,i} p_{G,i}$ s.t.  $p_{G(i)} - p_{D(i)} = \sum_{j \in \delta i} b_{ij} (\theta_i - \theta_j)$  ,  $p_{G,g}^{min} \leq p_{G,g} \leq p_{G,g}^{max}$ ,  $-f_{ij}^{max} \leq b_{ij} (\theta_i - \theta_j) \leq f_{ij}^{max},$  $i \in N, g \in \mathcal{G}, ij \in \mathcal{L}$ 

Benefit: We only look for active sets that occur in practice!

Learning optimal active sets through sampling

 $\begin{array}{cccc} \mathsf{load} & \mathsf{Optimal} & p_G^* & \mathsf{Optimal dispatch} \\ \mathsf{profile} \ \mathbf{i} & \longrightarrow & \mathsf{Power} & \longrightarrow & \mathcal{A}^* \\ \mathbf{p}_D & \mathsf{Flow} & \mathcal{A}^* & \mathsf{Optimal active set } \mathbf{i} \end{array}$ 

When to stop? Combined **probability** of observed patterns  $\mathcal{A}^* = \{\mathcal{A}^{(1)}, \mathcal{A}^{(2)}, ...\}$  is high. 
$$\begin{split} \min_{p_{G},\theta} & \sum_{i \in \mathcal{G}} c_{1,i} p_{G,i} \\ \text{s.t.} & p_{G(i)} - \boldsymbol{p}_{\boldsymbol{D}(i)} = \sum_{j \in \delta i} b_{ij} (\theta_{i} - \theta_{j}) \,, \\ & p_{G,g}^{\min} \leq p_{G,g} \leq p_{G,g}^{\max} \,, \\ & -f_{ij}^{\max} \leq b_{ij} (\theta_{i} - \theta_{j}) \leq f_{ij}^{\max} \,, \\ & i \in N, g \in \mathcal{G} \,, ij \in \mathcal{L} \end{split}$$

Learning optimal active sets through sampling

 $\begin{array}{cccc} \mathsf{load} & \mathsf{Optimal} & & \\ \mathsf{profile} \ \mathbf{i} & \longrightarrow & \mathsf{Power} & \longrightarrow & \\ \mathbf{p}_{\mathcal{D}} & & \mathsf{Flow} & & & \\ \end{array} \xrightarrow{p_{\mathcal{G}}^*} \mathsf{Optimal \ dispatch} \\ \mathcal{A}^* \ \mathsf{Optimal \ active \ set} \ \mathbf{i} \end{array}$ 

When to stop? Combined **probability** of observed patterns  $\mathcal{A}^* = \{\mathcal{A}^{(1)}, \mathcal{A}^{(2)}, ...\}$  is high.

Combined probability of observed patterns  $\mathcal{A}^*$  = the probability that  $\mathcal{A}^*$  contains the optimal policy and give us an optimal solution for a new sample  $p_D$  
$$\begin{split} \min_{p_{G},\theta} & \sum_{i \in \mathcal{G}} c_{1,i} p_{G,i} \\ \text{s.t.} & p_{G(i)} - \boldsymbol{p}_{\boldsymbol{D}(i)} = \sum_{j \in \delta i} b_{ij} (\theta_{i} - \theta_{j}) \,, \\ & p_{G,g}^{min} \leq p_{G,g} \leq p_{G,g}^{max} \,, \\ & -f_{ij}^{max} \leq b_{ij} (\theta_{i} - \theta_{j}) \leq f_{ij}^{max} \,, \\ & i \in N, g \in \mathcal{G} \,, ij \in \mathcal{L} \end{split}$$



Learning optimal active sets through sampling

 $\begin{array}{cccc} \mathsf{load} & & \mathsf{Optimal} & \\ \mathsf{profile} \ \mathbf{i} & \longrightarrow & \mathsf{Power} & \longrightarrow & \\ p_{\mathcal{D}} & & \mathsf{Flow} & & \mathcal{A}^* \ \mathsf{Optimal} \ \mathsf{active} \ \mathsf{set} \ \mathbf{i} \end{array}$ 

When to stop? Combined **probability** of observed patterns  $\mathcal{A}^*$  is high.

Requires estimating this probability!

	$\min_{p_G,\theta} \sum_{i\in \mathcal{G}} c_{1,i} p_{G,i}$
7	s.t.
	$p_{G(i)} - p_{D(i)} = \sum_{j \in \delta i} b_{ij} (\theta_i - \theta_j)$ ,
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T	$-f_{ij}^{max} \le b_{ij} (\theta_i - \theta_j) \le f_{ij}^{max},$
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Learning optimal active sets through sampling

 $\begin{array}{cccc} \mathsf{load} & \mathsf{Optimal} & p_G^* & \mathsf{Optimal dispatch} \\ \mathsf{profile} \ \mathbf{i} & \longrightarrow & \mathsf{Power} & \longrightarrow & \mathcal{A}^* & \mathsf{Optimal active set} \ \mathbf{i} \\ p_D & \mathsf{Flow} & \mathcal{A}^* & \mathsf{Optimal active set} \ \mathbf{i} \end{array}$ 

When to stop? Combined **probability** of observed patterns  $\mathcal{A}^*$  is high.

Requires estimating this probability!

Algorithm based on a "stream" of samples: [Misra, Roald and Ng '18]

 $\min_{p_G,\theta} \sum_{i\in\mathcal{G}} c_{1,i} p_{G,i}$ s.t.  $p_{G(i)} - p_{D(i)} = \sum_{j \in \delta i} b_{ij} (\theta_i - \theta_j)$  ,  $p_{G,g}^{min} \leq p_{G,g} \leq p_{G,g}^{max}$ ,  $-f_{ij}^{max} \le b_{ij} (\theta_i - \theta_j) \le f_{ij}^{max},$  $i \in N, g \in \mathcal{G}$  ,  $ij \in \mathcal{L}$ 

Learning optimal active sets through sampling

 $\begin{array}{cccc} \mathsf{load} & \mathsf{Optimal} & p_G^* & \mathsf{Optimal dispatch} \\ \mathsf{profile} \ \mathbf{i} & \longrightarrow & \mathsf{Power} & \longrightarrow & \mathcal{A}^* \\ p_D & \mathsf{Flow} & \mathcal{A}^* & \mathsf{Optimal active set } \mathbf{i} \end{array}$ 

When to stop? Combined **probability** of observed patterns  $\mathcal{A}^*$  is high.

Requires estimating this probability!

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- 1. Use first samples to identify optimal active sets  $\mathcal{A}^*$
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$$\begin{split} \min_{p_{G},\theta} & \sum_{i \in \mathcal{G}} c_{1,i} p_{G,i} \\ \text{s.t.} & p_{G(i)} - \boldsymbol{p}_{\boldsymbol{D}(i)} = \sum_{j \in \delta i} b_{ij} (\theta_{i} - \theta_{j}), \\ & p_{G,g}^{min} \leq p_{G,g} \leq p_{G,g}^{max}, \\ & -f_{ij}^{max} \leq b_{ij} (\theta_{i} - \theta_{j}) \leq f_{ij}^{max}, \\ & i \in N, g \in \mathcal{G}, ij \in \mathcal{L} \end{split}$$

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- 2. Use next samples to assess the "rate of discovery" (i.e., empirical probability of observing previously unseen  $\mathcal{A}^*$ )
- 3. Use empirically observed probability to derive bounds on the true probability
- 4. If true probability is too high: Add additional samples and move to step 1

 $\begin{array}{l} \min_{p_{G},\theta} \quad \sum_{i \in \mathcal{G}} c_{1,i} p_{G,i} \\
\text{s.t.} \\
p_{G(i)} - \boldsymbol{p_{D(i)}} = \sum_{j \in \delta i} b_{ij} (\theta_{i} - \theta_{j}), \\
p_{G,g}^{min} \leq p_{G,g} \leq p_{G,g}^{max}, \\
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Learning optimal active sets through sampling

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When to stop? Combined **probability** of observed patterns  $\mathcal{A}^*$  is high.

Requires estimating this probability!

Algorithm based on a "stream" of samples: [Misra, Roald and Ng '18]

- Guaranteed to converge
- Converges fast if there are few active sets
- No assumptions on problem structure/probability distribution required!



Learning optimal active sets through sampling

 $\begin{array}{cccc} \mathsf{load} & \mathsf{Optimal} & & p_G^* & \mathsf{Optimal dispatch} \\ \mathsf{profile} \ \mathbf{i} & \longrightarrow & \mathsf{Power} & \longrightarrow & \\ \mathbf{p}_D & & \mathsf{Flow} & & \mathcal{A}^* & \mathsf{Optimal active set } \mathbf{i} \end{array}$ 

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Algorithm based on a "stream" of samples: [Misra, Roald and Ng '18]





Missing mass problem and Good-Turing estimator! [Bertsimas and Stellato '18]

#### Using the Patterns

Operating the system based on learned policies



	$\min_{p_G,\theta} \sum_{i\in\mathcal{G}} c_{1,i} p_{G,i}$	
A	s.t. $p_{G(i)} - \mathbf{p}_{\mathbf{D}(i)} = \sum_{j \in \delta i} b_{ij} (\theta_i - \theta_j),$	
	$p_{G,g}^{min} \leq p_{G,g} \leq p_{G,g}^{max}$ ,	/
T	$-f_{ij}^{max} \le b_{ij} (\theta_i - \theta_j) \le f_{ij}^{max},$	
	$i \in N, g \in \mathcal{G}, ij \in \mathcal{L}$	



Operator experience!

Classification [Deka and Misra, 2019]

#### Using the Patterns

Operating the system based on learned policies

Load  $p_D$ 

$$\begin{split} \min_{p_{G},\theta} & \sum_{i \in \mathcal{G}} c_{1,i} p_{G,i} \\ \text{s.t.} & p_{G(i)} - \boldsymbol{p_{D(i)}} = \sum_{j \in \delta i} b_{ij} (\theta_{i} - \theta_{j}), \\ & p_{G,g}^{min} \leq p_{G,g} \leq p_{G,g}^{max}, \\ & -f_{ij}^{max} \leq b_{ij} (\theta_{i} - \theta_{j}) \leq f_{ij}^{max}, \\ & i \in N, g \in \mathcal{G}, ij \in \mathcal{L} \end{split}$$

Operator experience / classification [Deka and Misra, 2019]

- Works well if the number of optimal active sets is small
- Dangerous if a correct active set has not yet been observed

 $\mathcal{A}^{(1)} \longrightarrow \mathcal{P}_{G}$  $\mathcal{A}^{(2)} \longrightarrow \mathcal{P}_{G}^{(2)}$ 

 $\mathcal{A}^{(n)} \longrightarrow p_{c}^{(n)}$ 

#### Is it effective?

Stopping criterion: Rate of discovery (Empirical probability) < 0.01

- Works well if the number of optimal active sets is small
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#### Is it effective?

#### Stopping criterion: Rate of discovery (Empirical probability) < 0.01

		Normal distribution					Uniform distribu <mark>o</mark> n					
		$K_M$	M	$W_M$	R , $W$	$\mathbb{P}(p^*)$	$K_M$	M	$W_M$	$\overline{K}_{W}$	$\mathbb{P}(p^*)$	
	Low-Complexity											
	case3_1mbd	1	1	13'259	0.0	1.0	1	1	13'259	0.0	0.0	
	case5_pjm	1	1	13'259	0.0	1.0	1	1	13'259	0.0	0.0	
	case14_ieee	1	1	13'259	0.0	1.0	1	1	13'259	0.0	0.0	
	case30_ieee	1	1	13'259	0.0	1.0	1	1	13'259	0.0	0.0	
Tost	case39_epri	2	2	13'259	0.0	1.0	2	2	13'259	0.0008	0.9998	
1031	case118_ieee	2	33	13'259	0.0	1.0	2	4	13'259	0.0019	0.9984	
case	case57_ieee	2	2	13'259	0.0003	0.9997	3	46	13'259	0.0	0.0	
	case1888_rte	3	6	13'259	0.0	0.0	3	10	13'259	0.0	0.0	
	case1951_rte	5	47	13'259	0.0069	0.9943	11	63	13'259	0.0084	0.9901	
	case162_ieee_dtc	7	91	13'259	0.0054	0.9925	17	192	13'259	0.0085	0.9926	
	case24_ieee_rts	10	1456	18'209	0.0	0.0	11	64	13'259	0.0047	0.9941	
	High-Complexity											
	case73_ieee_rts	19	1258	17'844	0.0087	0.9931	130	22'000	24'977	0.0136	-	
	case300_ieee	24	1257	17'842	0.0073	0.9919	293	9095	22'789	0.0099	0.9897	
	case200_pserc	174	4649	21'112	0.0099	0.9909	236	6741	22'040	0.0099	0.9901	
	case240_pserc	2993	22'000	24'997	0.0795	-	2993	22'000	24'997	0.0795	-	

For many DC OPF test cases\* it works well:

Few active sets required to achieve a high probability

- Works well if the number of optimal active sets is small
- Dangerous if a correct active set has not yet been observed

#### Is it effective?

#### # of active sets

# of active sets

			Normal distribution					Uniform distribution					
			M	$W_M$	$R_{M,W}$	$\mathbb{P}(p^*)$	М	M	$W_M$	$R_{M,W}$	$\mathbb{P}(p^*)$		
	Low-Complexity												
	case3_lmbd	1	1	13'259	0.0	1.0	1	1	13'259	0.0	0.0		
	case5_pjm	1	1	13'259	0.0	1.0	1	1	13'259	0.0	0.0		
	case14_ieee	1	1	13'259	0.0	1.0	1	1	13'259	0.0	0.0		
	case30_ieee	1	1	13'259	0.0	1.0	1	1	13'259	0.0	0.0		
Test case	case39_epri	2	2	13'259	0.0	1.0	2	2	13'259	0.0008	0.9998		
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Few active sets required to achieve a high probability

- Works well if the number of optimal active sets is small
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#### **PSERC 200 bus test case**



Many optimal active sets with non-negligible probability = the algorithm terminates slowly!

#### **RTE 1951 bus test case**



#### Few optimal active sets with high probability = the algorithm terminates fast!

# How do the patterns change with renewable energy?



Renewable energy = Increasing load variability = Increasing number of active sets

Results for IEEE 300 bus system with normally distributed load uncertainty [Ng, Misra, Roald, Backhaus, 2018]



Renewable energy = Increasing load variability = Increasing number of active sets

Results for IEEE 300 bus system with normally distributed load uncertainty [Ng, Misra, Roald, Backhaus, 2018]

"Operators sometimes make wrong decisions because of increased variability"

Patrick Panciatici (RTE, France)

More variety in the active (congested) line constraints!

Without

Energy



Active line constraints for the month of June in IEEE RTS-GMLC test case (with and without renewable energy).

Red and blue:

More variety in the active (congested) line constraints!



Active line constraints for the month of June in IEEE RTS-GMLC test case (with and without renewable energy).

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Active line constraints for the month of June in IEEE RTS-GMLC test case (with and without renewable energy).

Red and blue:

Red and blue:

More variety in the active (congested) line constraints!



With more variable power injections, it gets harder to learn all the relevant patterns!

## How about AC optimal power flow?

#### Learning AC congestion patterns

iteration



AC OPF is **typically** harder (voltage and reactive power constraints)

but not always

#### Learning AC active constraints

#### **RTE 1951 bus test case**

#### **PSERC 200 bus test case**



If there are too many active sets, we can still learn active constraints.

#### Learning AC active constraints

**RTE 1951 bus test case** 

Only 164 of 5192 transmission line constraints ever active

#### **PSERC 200 bus test case**

Only 28 of 490 transmission line constraints ever active





If there are too many active sets, we can still learn active constraints.

## Only very few constraints seem to be at the optimal solution

## Are there constraints that can never be active at all?

#### Learning the active constraints

Previous part of the talk: If you know the set of active constraints at optimum, you can recover the optimal solution!



- Optimal active set = "minimal" information we need to recover optimal solution
- Inherently encodes information about physical constraints and technical limits
- Finite, low dimensional object
- Nice practical interpretation as "operational patterns"

#### Learning the active constraints

If you know the *set of active constraints* at optimum, you can recover the optimal solution!

If you know that some constraints are *never active*, you can solve a smaller optimization problem!

Previously, we used learning to identify constraints that are probably not active.

Let's see if we can find transmission constraints that will never be active (i.e., are redundant).

#### How many constraints can ever be active?

Previous part of the talk: If you know the set of active constraints at optimum, you can recover the optimal solution!

If you know that some constraints are *never active*, you can solve a solve a smaller optimization problem!



## 1) Parallel lines

- For **parallel lines**, one flow limit may always be more restrictive.
- Evaluating analytic condition is sufficient to identifying redundant flow limits on parallel lines using only the lines' parameters.

$$\frac{b_{ij}^{(k)}}{f_{ij}^{max\,(k)}} < \frac{b_{ij}^{(l)}}{f_{ij}^{max\,(l)}}$$

Case	Num.	Num. Parallel	Num. Redundant
Name	Lines	Lines	Limits
PL-2383wp	2896	20	6
PL-2736sp	3269	12	4
PL-2737sop	3269	12	4
PL-2746wop	3307	16	5
PL-2746wp	3279	12	4
PL-3012wp	3572	12	5
PL-3120sp	3693	18	8
PL-3375wp	4161	178	5
PEGASE-89	210	8	2
PEGASE-1354	1991	519	203
PEGASE-2869	4582	1157	316
PEGASE-9241	16049	3503	650

[Molzahn '18]

• For some test cases, **approximately 20% of parallel lines** are identified as redundant.

#### 2) Optimization-based constraint screening

min/max<br/> $p_G, p_D$  $f_l = M_{(l,\cdot)}(p_G - p_D)$ Minimize/maximize line flowss.t. $\sum_{i=1}^{N_B} (p_{G(i)} - p_{D(i)}) = 0$ Power balanceNon-redundant $\mathbf{0} \le p_G \le p_G^{max}$ ,Generation constraintsTypically non-redundant $-f_L^{max} \le M(p_G - p_D) \le f_L^{max}$ ,Transmission constraintsOften redundant $(1 - X) p_D^{nom} \le p_D \le (1 + X) p_D^{nom}$ , Considered load variationOften redundant

Allow power demand  $p_D$  to vary  $\pm X \cdot 100\%$  where  $0 \le X \le 1$ 

**Relax** generator lower bounds to **0** (applicable to unit commitment!)

Find maximum and minimum achievable flows  $\overline{f_l}$ ,  $f_l$
# 2) Optimization-based constraint screening

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Allow power demand  $p_D$  to vary  $\pm X \cdot 100\%$  where  $0 \le X \le 1$ 

**Relax** generator lower bounds to **0** (applicable to unit commitment!)

Find maximum and minimum achievable flows  $\overline{f_l}$ ,  $f_l$ 

If  $\overline{f_l} < f_{L,l}^{\max}$  or  $f_l > f_{L,l}^{\min} \Rightarrow$  Constraint is redundant!

#### Many redundant constraints...

Percentage of line flow constraints





Constraint screening results for test cases from PGLib v.17.08

## ...even for large load variations!

## Significant reduction in computation...



Single-period Unit Commitment problem

Results based on 100 computations of UC for test cases from PGLib v.17.08

...even for large load variations!

# Summary

Congestion pattern = Optimal active set of OPF

In traditional operations, only a few constraints are relevant

... our algorithms can identify them!

Renewable energy variability increases complexity of system operations = increasing number of operational patterns!

A good case for optimization?  $\ensuremath{\textcircled{\sc 0}}$ 



### Thank you!

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Yee Sian Ng, Sidhant Misra, Line Roald and Scott Backhaus, «Statistical Learning for DC Optimal Power Flow», Power System Computation Conference (PSCC), 2018

Sidhant Misra, Line Roald and Yee Sian Ng, «Learning for Constrained Optimization», https://arxiv.org/abs/1802.09639

L.A. Roald and D.K. Molzahn, "Implied Constraint Satisfaction in Power System Optimization: The Impacts of Load Variations," *57th Annual Allerton Conference on Communication, Control, and Computing,* 2019.